# E201B: Midterm Exam—Suggested Answers 

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## Risk Aversion

We have the utility function $u(c)=-e^{-c}$. The coefficients of absolute risk aversion, $A(c)$, and relative risk aversion, $R(c)$ are obtained from the following formulae:

$$
\begin{equation*}
A(c)=-\frac{u^{\prime \prime}(c)}{u^{\prime}(c)}, \quad R(c)=-\frac{u^{\prime \prime}(c) c}{u^{\prime}(c)} \tag{1}
\end{equation*}
$$

For this particular choice of utility, the coefficients are:

$$
\begin{equation*}
A(c)=-\frac{-e^{-c}}{e^{-c}}=1, \quad R(c)=-\frac{-e^{-c} c}{e^{-c}}=c \tag{2}
\end{equation*}
$$

Taking derivatives, it follows that $A^{\prime}(c)=0$, showing that absolute risk aversion is constant, and that $R^{\prime}(c)=1$, showing that relative risk aversion is increasing.

[^0]We are asked to consider the following game. (Bold-face numbers denote best responses.)

| $H$ | $G$ |  |
| :---: | :---: | :---: |
| $H$ | 0,0 | $\mathbf{2 , 1}$ |
| $G$ | $\mathbf{1 , 2}$ | 0,0 |
|  |  |  |

There are two pure-strategy Nash equilibria to this game: $(H, G)$ and $(G, H)$, with respective payoffs $(2,1)$ and $(1,2)$. Neither of these Nash equilibria is symmetric. There is also a mixedstrategy Nash equilibrium. It can be computed as follows. Assuming that both players are mixing between $H$ and $G$, it follows that they must each be indifferent between either strategy. Hence, if player II's probability of playing $H$ is $\beta$, then

$$
\begin{equation*}
E_{\beta}\left[u_{I}(H)\right]=E_{\beta}\left[u_{I}(G)\right] \quad \Rightarrow \quad 0 \beta+2(1-\beta)=\beta+0(1-\beta) \quad \Rightarrow \quad \beta=2 / 3 . \tag{3}
\end{equation*}
$$

By symmetry of the payoff bimatrix, it follows that $\alpha$, player I's probability of playing $H$, is also $2 / 3$. Hence, the mixed-strategy Nash equilibrium $(2[H] / 3+[G] / 3,2[H] / 3+[G] / 3)$ is symmetric. The vector of payoffs associated with this equilibrium is $(2 / 3,2 / 3)$.

The following correlated strategy (call it $\mu$ ) is a symmetric correlated equilibrium that Paretodominates the symmetric Nash.

|  | $H$ | $G$ |
| :---: | :---: | :---: |
|  |  | 0 |
|  | $1 / 2$ |  |
|  | $1 / 2$ | 0 |
|  |  |  |

Clearly this is symmetric. The vector of payoffs associated with this strategy is $(1.5,1.5)$, which clearly Pareto-dominates the symmetric Nash payoffs. To show that $\mu$ is indeed a correlated equilibrium, first notice that - by symmetry - if the incentive constraints are satisfied for player I then they are also satisfied for player II, so it suffices to check the constraints for player I only. If the mediator recommends player I to play $H$ then player I will obey the recommendation if

$$
\begin{equation*}
E_{\mu}\left[u_{I}(H) \mid H\right] \geq E_{\mu}\left[u_{I}(G) \mid H\right] . \tag{4}
\end{equation*}
$$

If the mediator recommends player I to play $H$, then player I knows that the mediator must have recommended player II to play $G$, i.e., player I's posterior probability over his opponent's actions given $H$ is that $G$ will be played with probability 1 . Hence, $E_{\mu}\left[u_{I}(H) \mid H\right]=2$, which is clearly greater than $0=E_{\mu}\left[u_{I}(G) \mid H\right]$.

If the mediator recommends player I to play $G$ then player I will obey the recommendation if

$$
\begin{equation*}
E_{\mu}\left[u_{I}(G) \mid G\right] \geq E_{\mu}\left[u_{I}(H) \mid G\right] . \tag{5}
\end{equation*}
$$

If the mediator recommends player I to play $G$, then player I knows that the mediator must have recommended player II to play $H$, i.e., player I's posterior probability over his opponent's actions given $G$ is that $H$ will be played with probability 1 . Hence, $E_{\mu}\left[u_{I}(G) \mid G\right]=1$, which is clearly greater than $0=E_{\mu}\left[u_{I}(H) \mid G\right]$. It now follows that $\mu$ is a correlated equilibrium.

## Trembling Hand Perfection

Definition $1 A$ strategy profile $\sigma$ is trembling hand perfect if there exists a sequence of strategy profiles $\left\{\sigma^{n}\right\}$ with $\sigma^{n} \rightarrow \sigma$ such that $\sigma_{i}^{n}\left(a_{i}\right)>0$ for every $a_{i} \in A_{i}$ and $\sigma_{i}\left(a_{i}\right)>0$ implies that $a_{i} \in B_{i}\left(\sigma_{-i}^{n}\right)$, where $B_{i}(\cdot)$ is player $i$ 's best-response correspondence.

Claim 2 Every trembling hand perfect equilibrium is a Nash equilibrium.

Proof-Let $\sigma$ be a trembling hand perfect equilibrium. Then there is a sequence $\left\{\sigma^{n}\right\}$ that satisfies the conditions of Definition 1. Fix any player $i$, and fix any action $a_{i} \in A_{i}$ such that $\sigma_{i}\left(a_{i}\right)>0$. By hypothesis,

$$
\begin{equation*}
u_{i}\left(a_{i}, \sigma_{-i}^{n}\right) \geq u_{i}\left(b_{i}, \sigma_{-i}^{n}\right), \quad \forall b_{i} \in A_{i} . \tag{6}
\end{equation*}
$$

Since $u_{i}$ is continuous, $\sigma_{-i}^{n} \rightarrow \sigma_{-i}$ implies that $u_{i}\left(\cdot, \sigma_{-i}^{n}\right) \rightarrow u_{i}\left(\cdot, \sigma_{-i}\right)$. Therefore, taking limits on each side of (6), we may conclude that

$$
\begin{equation*}
u_{i}\left(a_{i}, \sigma_{-i}\right) \geq u_{i}\left(b_{i}, \sigma_{-i}\right), \quad \forall b_{i} \in A_{i} . \tag{7}
\end{equation*}
$$

Notice finally that this condition applies to any action $a_{i}$ with $\sigma_{i}\left(a_{i}\right)>0$. Therefore for any two $a_{i}$ and $a_{i}^{\prime}$ such that $\sigma_{i}\left(a_{i}\right)>0$ and $\sigma\left(a_{i}^{\prime}\right)>0$, we must have both $u_{i}\left(a_{i}, \sigma_{-i}\right) \geq u_{i}\left(a_{i}^{\prime}, \sigma_{-i}\right)$ and $u_{i}\left(a_{i}^{\prime}, \sigma_{-i}\right) \geq u_{i}\left(a_{i}, \sigma_{-i}\right)$, so $u_{i}\left(a_{i}, \sigma_{-i}\right)=u_{i}\left(a_{i}^{\prime}, \sigma_{-i}\right)$. This, together with (7), implies that $\sigma$ be a Nash equilibrium.

## Subgame Perfection

To find a subgame perfect equilibrium we must show that it's Nash in every subgame. We will find it by backward induction. In the subgame below, there is only one Nash: $(U, L)$.

|  | $L$ | $R$ |
| :---: | :---: | :---: |
| $U$ | 6,6 | 3, 0 |
| $D$ | 0,0 | 2, 2 |

This is apparent from the fact that $D$ is strictly dominated by $U$.
Proceeding up the game tree, player I has the option of ending the game (playing $E$ ) or participating in the subgame above (playing $P$ ). Payoffs for each player assuming that the unique Nash equilibrium will be played in the subgame above are as follows.

$$
\begin{array}{l|l|}
\hline & 4,4 \\
\cline { 2 - 3 } & \mathbf{6 , 6} \\
\hline
\end{array}
$$

Therefore the unique subgame perfect equilibrium is $((P, U), L)$.

The strategic-form representation of this game looks like this.

|  | $L$ | $R$ |
| :---: | :---: | :---: |
| $(P, U)$ | $\mathbf{6 , 6}$ | 3,0 |
| $(P, D)$ | 0,0 | $2, \mathbf{2}$ |
| $(E, U)$ | $4, \mathbf{4}$ | $\mathbf{4 , 4}$ |
| $(E, D)$ | $4, \mathbf{4}$ | $\mathbf{4} \mathbf{4}$ |
|  |  |  |

Applying iterated weak dominance, we eliminate $(P, D)$ since it's strictly dominated by $(P, U)$. But then $R$ is weakly dominated by $L$, and finally, restricted to $L,(P, U)$ dominates the rest of player I's remaining strategies, leaving only one profile surviving, $((P, U), L)$. This is the subgame perfect equilibrium profile.


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