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## **The Competitive Mechanism**

Cobb-Douglas Aggregate Excess Demans:

$$z_1 = -\frac{1}{2} + \frac{1}{2} \frac{p_2}{p_1}$$
$$z_2 = \frac{1}{2} \frac{p_1}{p_2} - \frac{1}{2}$$

Aggregate excess demand: two key properties

### Homogeneous of degree zero

$$z_j(\lambda p) = z_j(p)$$
$$z_1 = -\frac{1}{2} + \frac{1}{2}\frac{p_2}{p_1}$$

#### Walras's Law

$$\sum_{j=1}^{k} p_j z_j(p) = 0$$

$$p_1\left[-\frac{1}{2} + \frac{1}{2}\frac{p_2}{p_1}\right] + p_2\left[\frac{1}{2}\frac{p_1}{p_2} - \frac{1}{2}\right] = ??$$

### Solving for Equilibrium

there are k different excess demand conditions  $z_j(p) = 0$  and there are k different prices  $p_1, \ldots, p_k$ 

but one excess demand condition is redundant

suppose  $z_j(p) = 0$  for j = 1, ..., k - 1, then from Walras's law  $z_k(p) = 0$ 

on the other hand, if  $z_j(p) = 0$  for all j = 1,...,k then so does  $z_j(\lambda p) = 0$ 

so many competitive equilibria

can solve only for relative prices using k - 1 excess demand equations Existence of competitive equilibrium? The Numeraire

may arbitrarily set the price of one good to 1

called the numeraire good, all prices are measured relative to that good

(for example – money is numeraire)

### **Cobb-Douglas Example**

Pick one equation

$$z_1 = -\frac{1}{2} + \frac{1}{2}\frac{p_2}{p_1} = 0$$

so 
$$p_2 / p_1 = 1$$

pick the other equation

$$z_2 = \frac{1}{2} \frac{p_1}{p_2} - \frac{1}{2}$$

get the same answer of course

if we choose good 1 as numeraire then we have  $p_1 = 1, p_2 = 1$ 

how do we find individual demands?

### **The First Welfare Theorem**

Suppose we have a competitive equilibrium with prices p and individual demands  $x_i^j$ 

is this pareto efficient?

That is: can we find  $\tilde{x}_j^i$  socially feasible that makes nobody worse off and at least one person better off?

That is: can we find  $\sum_{i=1}^{n} \tilde{x}_{j}^{i} \leq \sum_{i=1}^{n} \overline{x}_{j}^{i}$  so that  $u^{i}(\tilde{x}^{i}) \geq u^{i}(\overline{x}^{i})$  for everybody (all *i*) and for somebody (some *i*)  $u^{i}(\tilde{x}^{i}) > u^{i}(\overline{x}^{i})$ ? Observation: if  $u^{i}(\tilde{x}^{i}) > u^{i}(\overline{x}^{i})$  then  $p \cdot \tilde{x}^{i} > p \cdot \overline{x}^{i}$ Why??

Further observation:  $u^i(\tilde{x}^i) \ge u^i(\overline{x}^i)$  then  $p \cdot \tilde{x}^i \ge p \cdot \overline{x}^i$  (otherwise spend your extra income to buy more)

Our conclusion: if  $u^i(\tilde{x}^i) \ge u^i(\overline{x}^i)$  for everybody (all *i*) and for somebody (some *i*)  $u^i(\tilde{x}^i) > u^i(\overline{x}^i)$ , then

 $p \cdot \tilde{x}^i \ge p \cdot \overline{x}^i$  for all i and  $p \cdot \tilde{x}^i > p \cdot \overline{x}^i$  for some i

add these together:

$$\sum_{i=1}^{n} p \cdot \tilde{x}^{i} > \sum_{i=1}^{n} p \cdot \overline{x}^{i}$$

on the other hand

 $\sum_{i=1}^{n} \tilde{x}_{j}^{i} \leq \sum_{i=1}^{n} \overline{x}_{j}^{i}, \text{ so adding over different goods}$  $\sum_{j=1}^{k} p_{j} \sum_{i=1}^{n} \tilde{x}_{j}^{i} \leq \sum_{j=1}^{k} p_{j} \sum_{i=1}^{n} \overline{x}_{j}^{i}$ which says that  $\sum_{i=1}^{n} p \cdot \tilde{x}^{i} \leq \sum_{i=1}^{n} p \cdot \overline{x}^{i}$ 

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- $\hfill\square$  relationship to the core
- Implications for international trade
- □ the edgeworth box
- the second welfare theorem
- the competitive mechanism

### **Finance**

Trade in period 0 claims to consumption in period 1

k different states of nature in period 1, probability of state j is  $\pi_i$ 

consumption in state j is  $c_j$ 

time 0 price of consumption in state j is  $p_j$ 

budget constraint  $\sum_{j=1}^{k} p_j c_j \leq m$ 

"martingale" prices  $\tilde{p}_j = p_j / \pi_j$ 

budget constraint in martingale prices

$$E\tilde{p}c = \sum_{j=1}^{k} \pi_j \tilde{p}_j c_j \le m$$

#### Securities

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Security a pays r_j in state j
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Examples:

Arrow security on state j pays 1 in state j 0 in all other states

Price of an arrow security  $p_i$  or  $\tilde{p}_j$ 

Bond pays 1 in all states

Price of a bond  $\sum_{j=1}^{k} p_j$  or  $E\tilde{p}$ 

"arbitrage pricing" = law of one price

### Spanning

Two states j = 1, 2, k = 2

Stock 1 pays (2,1) price  $q_1$ , Stock 2 pays (1,2) price  $q_2$ 

What is the price of a bond?

Buy both stocks: get (3,3) bond pays (1,1) so  $1/3^{rd}$  of both stocks

 $q_b = (q_1 + q_2)/3$ 

what does spanning mean?

k different assets that are "independent"

can determine all asset prices in terms of a spanning set of assets

#### Short Sales

Stock 1 pays (2,1) price  $q_1$ , Stock 2 pays (3,1) price  $q_2$ 

What is the price of a bond?

a(2,1) + b(3,1) = (1,1)  $\begin{bmatrix} 2 & 3 \\ 1 & 1 \end{bmatrix} \begin{bmatrix} a \\ b \end{bmatrix} = \begin{bmatrix} 1 \\ 1 \end{bmatrix}$   $\begin{bmatrix} a \\ b \end{bmatrix} = \begin{bmatrix} 2 & 3 \\ 1 & 1 \end{bmatrix}^{-1} \begin{bmatrix} 1 \\ 1 \end{bmatrix}$   $= (-1) \begin{bmatrix} 1 & -3 \\ -1 & 2 \end{bmatrix} \begin{bmatrix} 1 \\ 1 \end{bmatrix}$   $= \begin{bmatrix} 2 \\ -1 \end{bmatrix}$ 

check 2(2,1) - 1(3,1) = (4,2) - (3,1) = (1,1)

so buying 2 units of stock 1 and -1 units of stock 2 is the same as buying a bond: cost  $2q_1 - q_2$ 

what does it mean to buy -1 unit of stock 2?

observation  $2q_1 - q_2 > 0$  so we can conclude that  $q_1 > q_2 / 2$ 

# **The Holdup Problem**

"entrepreneur" (inventor, merchant) creates value of  $\rho$ 

 $\rho$  is drawn from a uniform distribution over [0,1] and is private information to the entrepreneur

case 1: the innovator receives a fraction of the social total  $\phi\rho$ 

case 2: the innovator receives the entire social total  $\rho$  but must pay N existing "rights holders" for the right to create value

examples:

the silk road

patents and copyrights

pollution

efficiency = the good is always produced

in case 1 the good is always produced

in case 2

rights holder *i* set price  $p_i$  for his right and gets an expected revenue of

$$\left(1-(N-1)p-p_i\right)p_i$$

p = 1/(N+1)

entrepreneur pays  $\frac{N}{N+1}$  to clear the needed rights, so creation if  $\frac{N}{N+1} < \rho$ 

what happens as  $N \to \infty$ 

as technologies grow more and more complex requiring more and more specialized inputs, monopoly power induced by patents and copyright becomes more and more socially damaging