

1. Principal Agent

a)

	0	1
Pay	0,0	1,1
Not Pay	0,0 [unique pure Nash]	5,-1

b) 1 can mix up to 50% on pay, 2 chooses 0

c) pay is weakly dominated for 1

d) pure is (pay,1); mixed is (50-50,1)

e) minmax for player 1 is 0; minmax for player 2 is 0

f) worst equilibrium is 0; best equilibrium is 1 provided

$$(1 - \delta)4 - \delta 1 \leq 0$$

$$4/5 \leq \delta$$

g) many answers: most obvious – play pay, 1 unless there has been a deviation, then switch to 0

h) convex hull of 0,0; 1,1; 5,-1 above the point 0,0

2. Auto Repair

a) normal form

	Repair	Not
Repair	$p - c, \theta v - p$	0, πv
Not	$p, \pi v - p$	0, πv [unique pure Nash]

b) some mixed Nash but all yield payoff 0; minmax for 1 is 0; pure precommitment is repair $p - c$; mixed precommitment mix so that 2 is indifferent

$$\alpha \theta v + (1 - \alpha) \pi v = \pi v + p$$

$$\alpha = \frac{p}{(\theta - \pi)v}$$

so payoff is $p - \alpha c$.

c) worst is obviously 0

d)

$$v = (1 - \delta)(p - c) + \delta(\theta v + (1 - \theta)w(n))$$

$$v = (1 - \delta)p + \delta(\pi v + (1 - \pi)w(n))$$

$$v = (1 - \delta)(p - c) + \delta(\theta v + (1 - \theta)w(n))$$

$$\frac{(1 - \delta\pi)}{(1 - \pi)}v - \frac{(1 - \delta)}{(1 - \pi)}p = \delta w(n)$$

$$v = (1 - \delta)(p - c) + \delta\theta v + (1 - \theta)\left[\frac{(1 - \delta\pi)}{(1 - \pi)}v - \frac{(1 - \delta)}{(1 - \pi)}p\right]$$

$$v\left[1 - \frac{1 - \theta + \delta\theta - \delta\pi}{(1 - \pi)}\right] = (1 - \delta)(p - c) - (1 - \theta)\frac{(1 - \delta)}{(1 - \pi)}p$$

$$v(\theta - \pi) = (1 - \pi)(p - c) - (1 - \theta)p = (\theta - \pi)p - (1 - \pi)c$$

$$v = p - \frac{1 - \pi}{\theta - \pi}c$$

$$\begin{aligned} w(n) &= \frac{(1 - \delta\pi)}{\delta(1 - \pi)}v - \frac{(1 - \delta)}{\delta(1 - \pi)}p \\ &= \frac{(1 - \delta\pi)}{\delta(1 - \pi)}\left[p - \frac{1 - \pi}{\theta - \pi}c\right] - \frac{(1 - \delta)}{\delta(1 - \pi)}p \\ &= p - \frac{(1 - \delta\pi)}{\delta(\theta - \pi)}c \geq 0 \end{aligned}$$

$$\delta \geq \frac{c}{p(\theta - \pi) + \pi c} < 1$$

$$(\theta - \pi)p - c > -\pi c$$

3. Auction

a) dominant strategy to announce truthfully; $\frac{3}{4}$ chance one person has low value and get revenue \$8; $\frac{1}{4}$ chance both have high value and get revenue \$12, so expected revenue is 9.

b) this is a 2x2 symmetric game; strategy is what to bid when type is \$12

calculate payoff to \$12 type only, since only that type has a choice

both bid \$12

$\frac{1}{2}$ opponent bids 12 and is 12: you win $\frac{1}{2}$ and get 0; lose $\frac{1}{2}$ and get 6, expected value 3

$\frac{1}{2}$ opponent bids 8 and is 8; you win always and get 4, expected value 4

you bid 8, he bids 12

$\frac{1}{2}$ opponent bids 12 and is 12; you lose and get 6; expected value 6

$\frac{1}{2}$ opponent bids 8 and is 8; you win $\frac{1}{2}$ and get 4; lose $\frac{1}{2}$ and get 4; expected value 4

you bid 12 he bids 8

you always win and get 4

both bid \$8

$\frac{1}{4}$ chance win versus 12, getting 4

$\frac{1}{4}$ chance lose versus 12, getting 6

$\frac{1}{4}$ chance win versus 8, getting 4

¼ chance lose versus 8, getting 4

	\$12	\$8
\$12	3.5,3.5	4,5
\$8	5,4	4.25, 4.25

So bidding 8 is dominant strategy equilibrium, and expected revenue to the seller is \$8.

4) Mechanism Design

$$\begin{aligned}
 & (1/3)[u(H, H; H) + u(L, H; L) + u(H, L; H)] \\
 & + (1/3)[t(H, H) + t(L, H) + t(H, L)] = \\
 & (x - 1)/3 + (1/3)[t(H, H) + t(L, H) + t(H, L)] \geq \\
 & (1/3)[u(H, H; H) + u(L, H; H) + u(H, L; H)] \\
 & + (1/3)[t(H, H) + t(H, H) + t(H, L)] = \\
 & x/3 + (1/3)[2t(H, H) + t(H, L)] \\
 & (x - 1)/3 + (1/3)[t(H, H) + t(L, H) + t(H, L)] \geq \\
 & x/3 + (1/3)[2t(H, H) + t(H, L)]
 \end{aligned}$$

$$t(L, H) - t(H, H) \geq 1$$

$t(H, H) = 0 \Rightarrow t(L, H) = 1$ and then choosing $t(H, L) = -1$ give the other constraint these constraints can be interpreted as “budget balance” when there are two players.